
ECE 307 – Techniques for Engineering Decisions

Lecture 2. Introduction to Linear Programming

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OUTLINE

- ❑ The nature of a programming or an optimization problem
- ❑ Linear programming (LP) : salient characteristics
- ❑ The LP problem formulation
- ❑ The LP problem solution
- ❑ Extensive illustrations with numerical examples

EXAMPLE 1: HIGH/LOW HEEL SHOE CHOICE PROBLEM

- ☐ **A lady is headed to a party and is trying to find a pair of shoes to wear; the choice is narrowed down to two possible choices:**
 - ☐ **a high heel pair; and**
 - ☐ **a low heel pair**
- ☐ **The high heel shoes look more beautiful but are not as comfortable as the competing pair**
- ☐ **Which pair should she choose?**

MODEL FORMULATION

- ❑ We first quantify our assessment along the two dimensions of *looks* and *comfort* in a table

<i>aspect</i>	<i>maximum value</i>	<i>assessment</i>		<i>weighting factor (%)</i>
		<i>high heels</i>	<i>low heels</i>	
<i>aesthetics</i>	5.0	4.2	3.6	70
<i>comfort</i>	5.0	3.5	4.8	30

- ❑ Next, we represent the decision in terms of two decision variables:

MODEL FORMULATION

$$x_H = \begin{cases} 1 & \text{choose high} \\ 0 & \text{otherwise} \end{cases} \quad x_L = \begin{cases} 1 & \text{choose low} \\ 0 & \text{otherwise} \end{cases}$$

- We formulate the objective to be the maximization of the *weighted* assessment

$$\max \{ 70 \% * \text{aesthetics} + 30 \% * \text{comfort} \}$$

- We state the objective in terms of the defined decision variables

$$\max Z = x_H [(4.2)(0.7) + (3.5)(0.3)] + x_L [(3.6)(0.7) + (4.8)(0.3)]$$

MODEL FORMULATION

□ Next, we consider the problem constraints:

○ only one pair of shoes can be selected

○ each decision variable is nonnegative

□ We express the constraints in terms of x_H and x_L

$$x_H + x_L = 1$$

$$x_H \geq 0, x_L \geq 0$$

PROBLEM STATEMENT SUMMARY

□ Decision variables:

$$x_H = \begin{cases} 1 & \text{choose high} \\ 0 & \text{otherwise} \end{cases} \quad x_L = \begin{cases} 1 & \text{choose low} \\ 0 & \text{otherwise} \end{cases}$$

□ Objective function:

$$\max Z = 3.99 x_H + 3.96 x_L$$

□ Constraints:

$$x_H + x_L = 1$$

$$x_H \geq 0, x_L \geq 0$$

THE OPTIMAL SOLUTION

- We determine the values x_H^* and x_L^* which result in the value of Z^* such that

$$Z^* = Z(x_H^*, x_L^*) \geq Z(x_H, x_L) \quad (\dagger)$$

for all *feasible* (x_H, x_L)

- We call such a solution *an optimal solution*
- A *feasible* solution is one that satisfies all the constraints
- The *optimal* solution, denoted by (x_H^*, x_L^*) , is selected from all the *feasible* solutions to the problem so as to satisfy (\dagger)

SOLUTION APPROACH: EXHAUSTIVE SEARCH

- We enumerate all the feasible solutions: in this problem there are only two alternatives:

$$A : \begin{cases} x_H = 1 \\ x_L = 0 \end{cases} \qquad B : \begin{cases} x_H = 0 \\ x_L = 1 \end{cases}$$

- We evaluate Z for A and B and compare

$$Z_A = 3.99$$

$$Z_B = 3.96$$

so that $Z_A > Z_B$ and so A is the optimal choice

- The *optimal* solution is

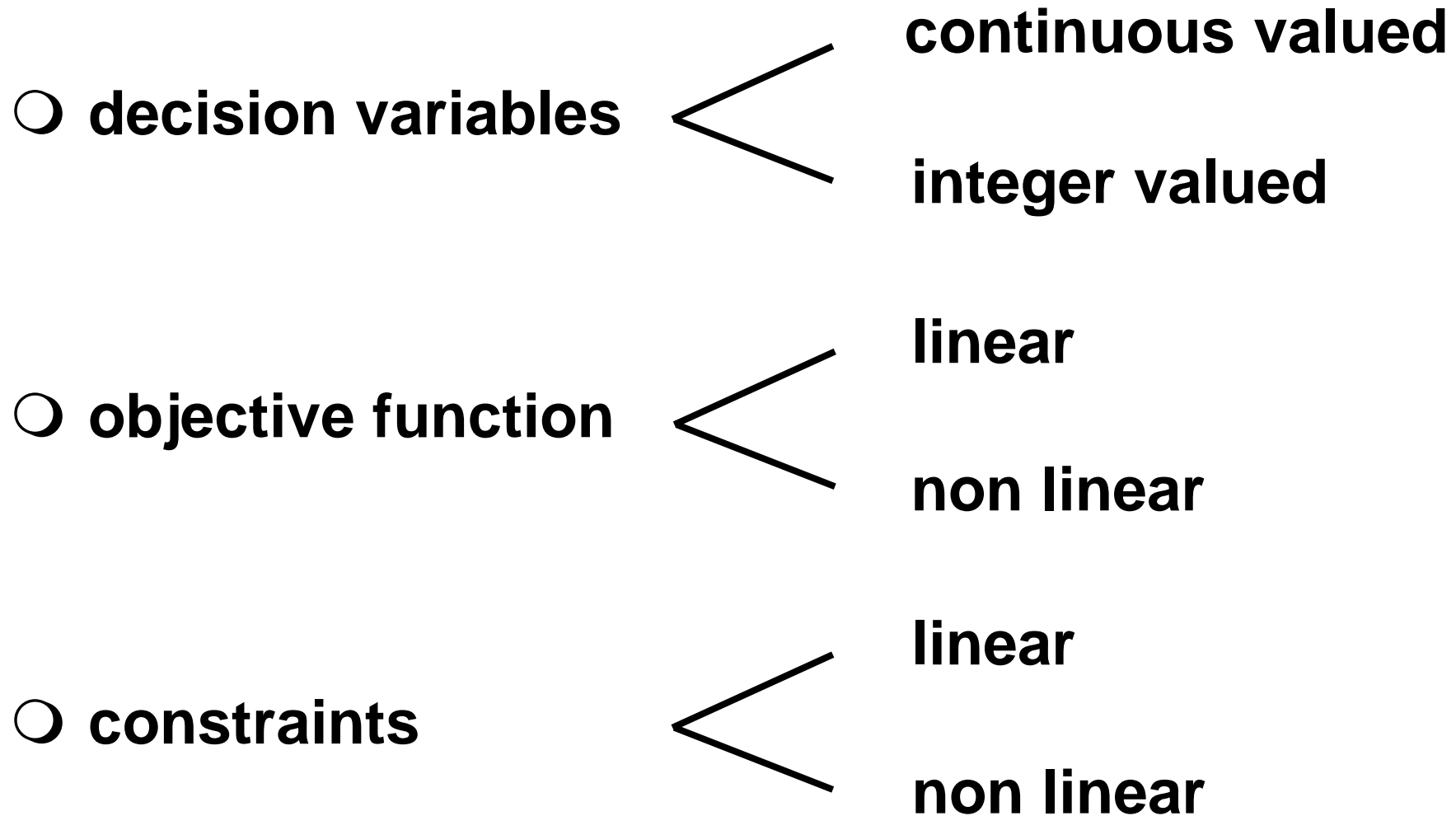
$$x_H^* = 1, \quad x_L^* = 0 \quad \text{and} \quad Z^* = 3.99$$

CHARACTERISTICS OF A PROGRAMMING/OPTIMIZATION PROBLEM

- ❑ The objective is to select the decision among the various alternatives and therefore requires first the *definition* of the *decision variables*
- ❑ We determine the “**best**” decision is on the basis of the objective function and so we need to obtain the *mathematical formulation* of the *objective function*
- ❑ The decision must satisfy *each specified constraint* and so we require the *mathematical statement* of the *problem constraints*

CLASSIFICATION OF PROGRAMMING PROBLEMS

The problem statement is characterized by :



PROGRAMMING PROBLEM CLASSES

- ☐ **Linear/nonlinear programming**
- ☐ **Static/dynamic programming**
- ☐ **Integer programming**
- ☐ **Mixed programming**

EXAMPLE 2: CONDUCTOR PROBLEM

- ❑ A company is producing two types of conductors for *EHV* transmission lines

<i>type</i>	<i>conductor</i>	<i>production capacity (unit/day)</i>	<i>metal needed (tons/unit)</i>	<i>profits (\$/unit)</i>
1	ACSR 84/19	4	1/6	3
2	ACSR 18/7	6	1/9	5

- ❑ The supply department can provide up to 1 *ton* of metal each day
- ❑ We schedule the production so as to *maximize* the profits of the company

PROBLEM ANALYSIS

- ❑ **Formulation of the objective: to *maximize* the profits of the company**
- ❑ **Means to attain this objective: determine how many units of product 1 and of product 2 to produce each day**
- ❑ **Consideration of all the constraints: the daily production capacity limits, the daily metal supply limit and *common sense* requirements**

MODEL CONSTRUCTION

□ We define the decision variables to be

x_1 = *number of type 1 units produced per day*

x_2 = *number of type 2 units produced per day*

□ We define the objective to be

Z = *profits (\$/day)*

$$= 3x_1 + 5x_2$$

□ *Sanity check* for units of the objective function

$$(\$/\text{day}) = (\$/\text{unit}) \cdot (\text{unit}/\text{day})$$

PROBLEM STATEMENT

❑ **Objective function:**

$$\max Z = 3x_1 + 5x_2$$

❑ **Constraints:**

○ **capacity limits:**

$$x_1 \leq 4 \quad x_2 \leq 6$$

○ **metal supply limit:**

$$\frac{x_1}{6} + \frac{x_2}{9} \leq 1$$

○ **common sense requirements:**

$$x_1 \geq 0, x_2 \geq 0$$

PROBLEM STATEMENT

$$\max Z = 3x_1 + 5x_2$$

s.t.

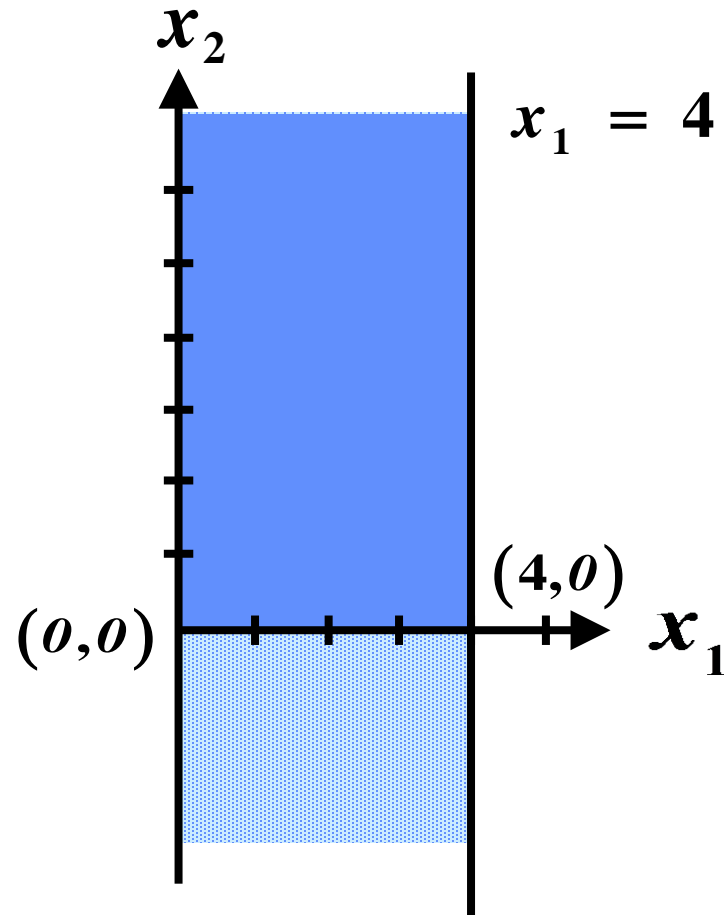
$$x_1 \leq 4$$

$$x_2 \leq 6$$

$$\frac{x_1}{6} + \frac{x_2}{9} \leq 1$$

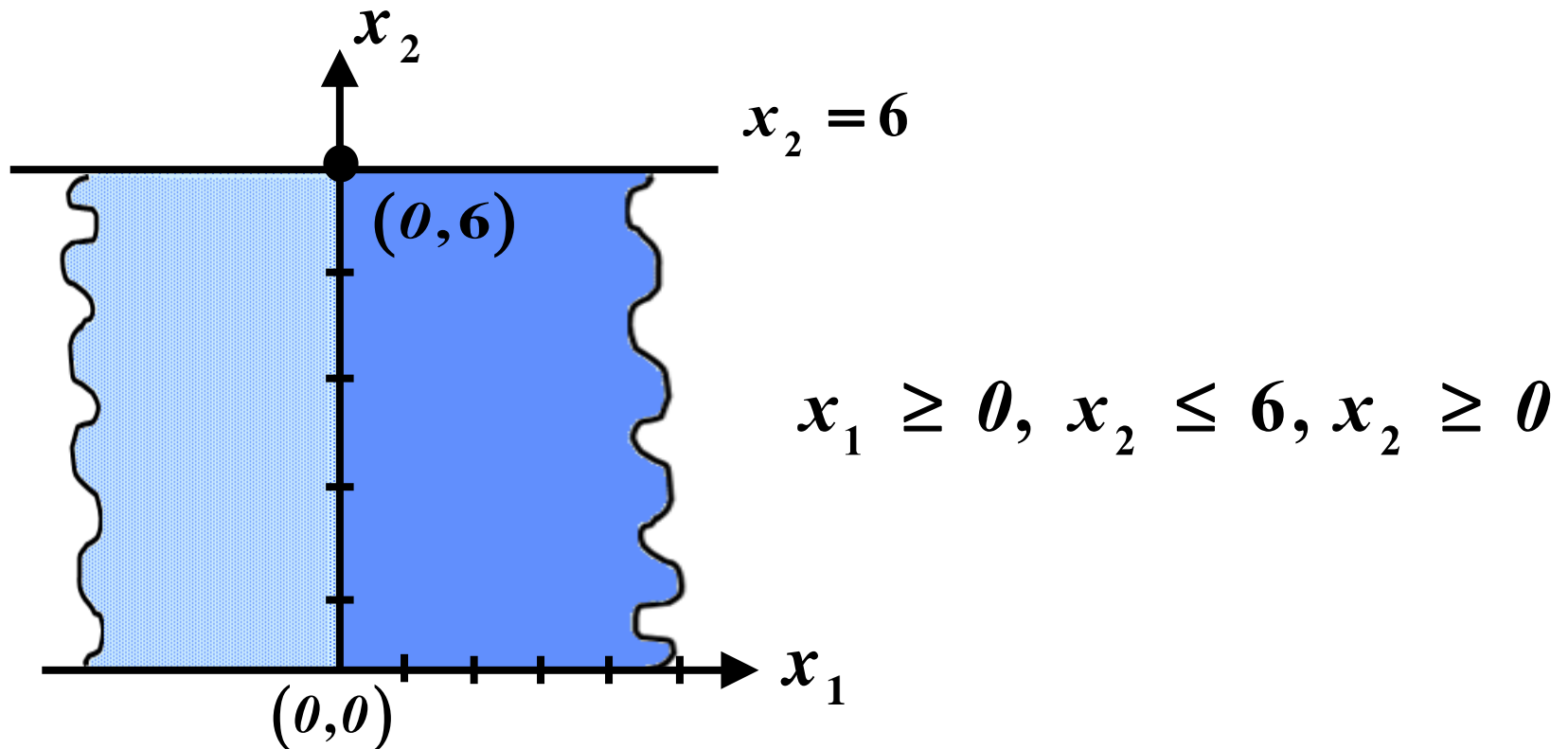
$$x_1 \geq 0, \quad x_2 \geq 0$$

VISUALIZATION OF THE *FEASIBLE REGION*

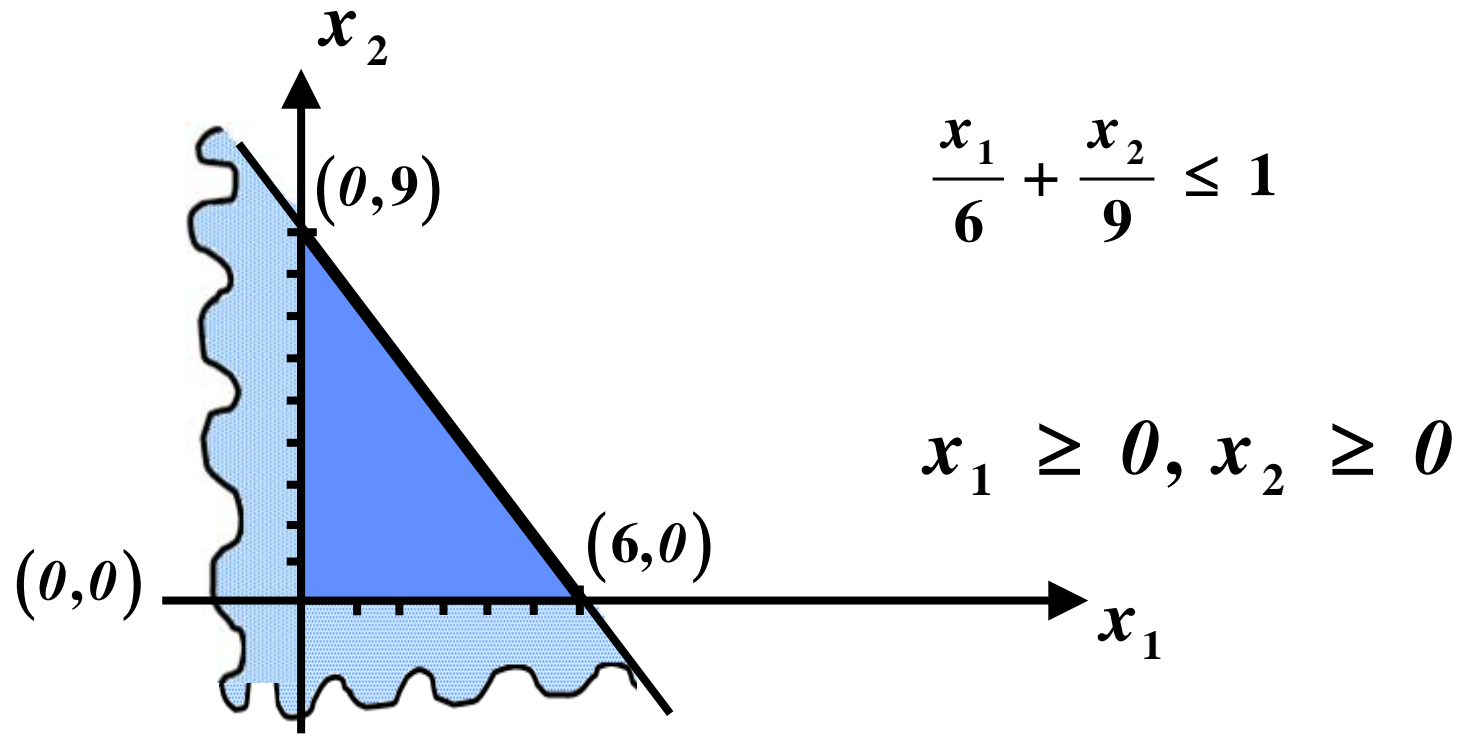


$$x_1 \geq 0, \quad x_1 \leq 4, \quad x_2 \geq 0$$

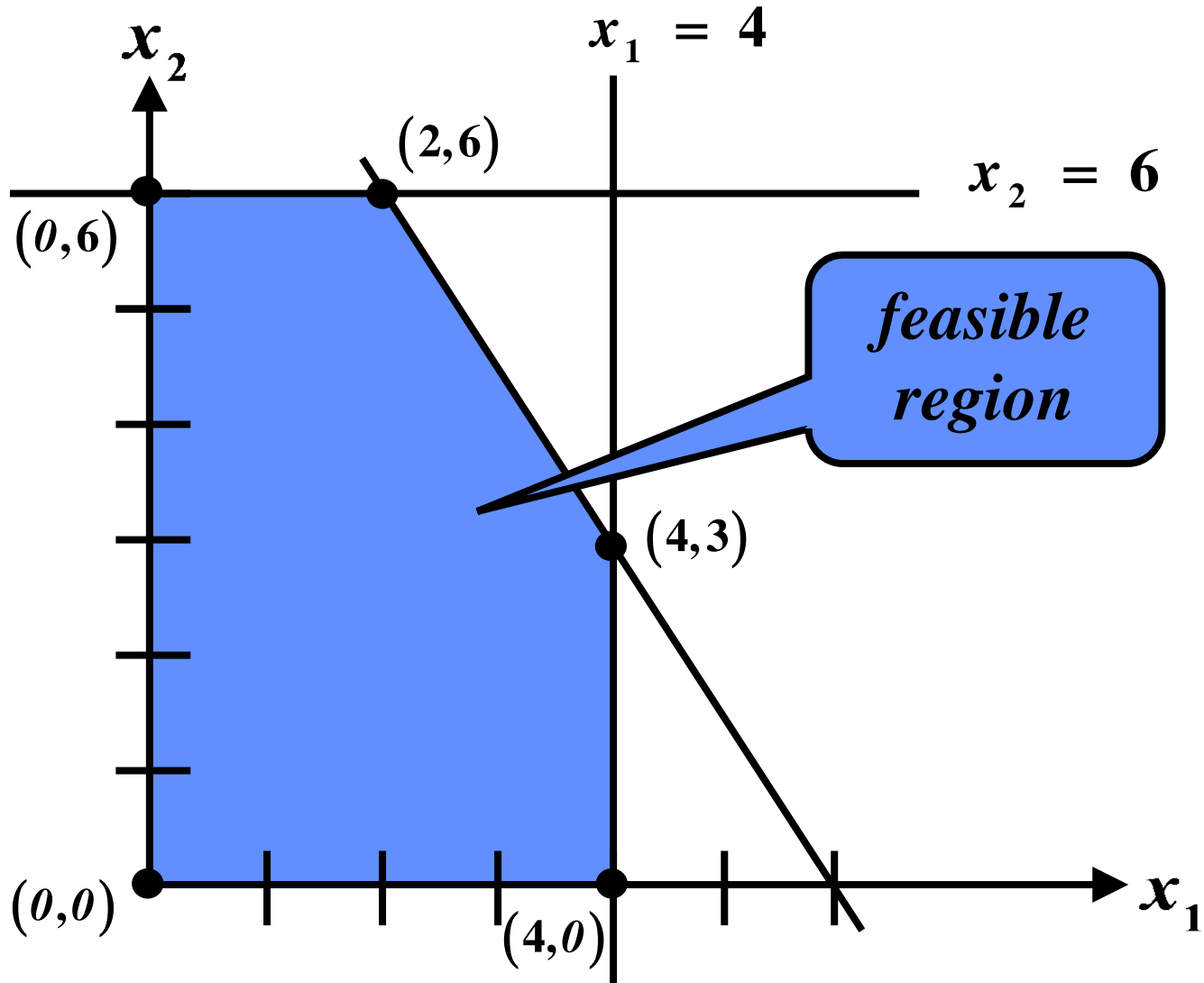
VISUALIZATION OF THE *FEASIBLE REGION*



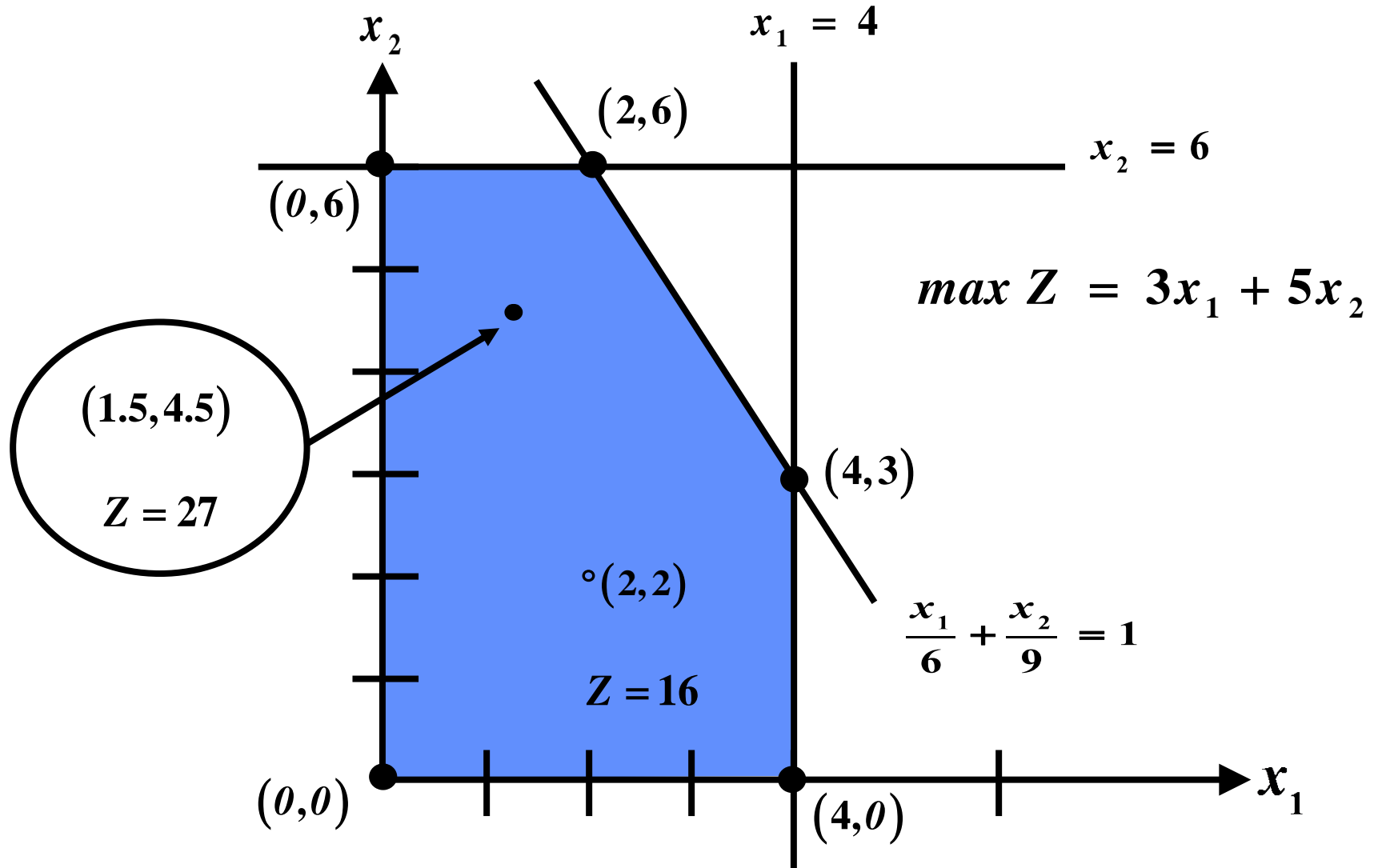
VISUALIZATION OF THE *FEASIBLE REGION*



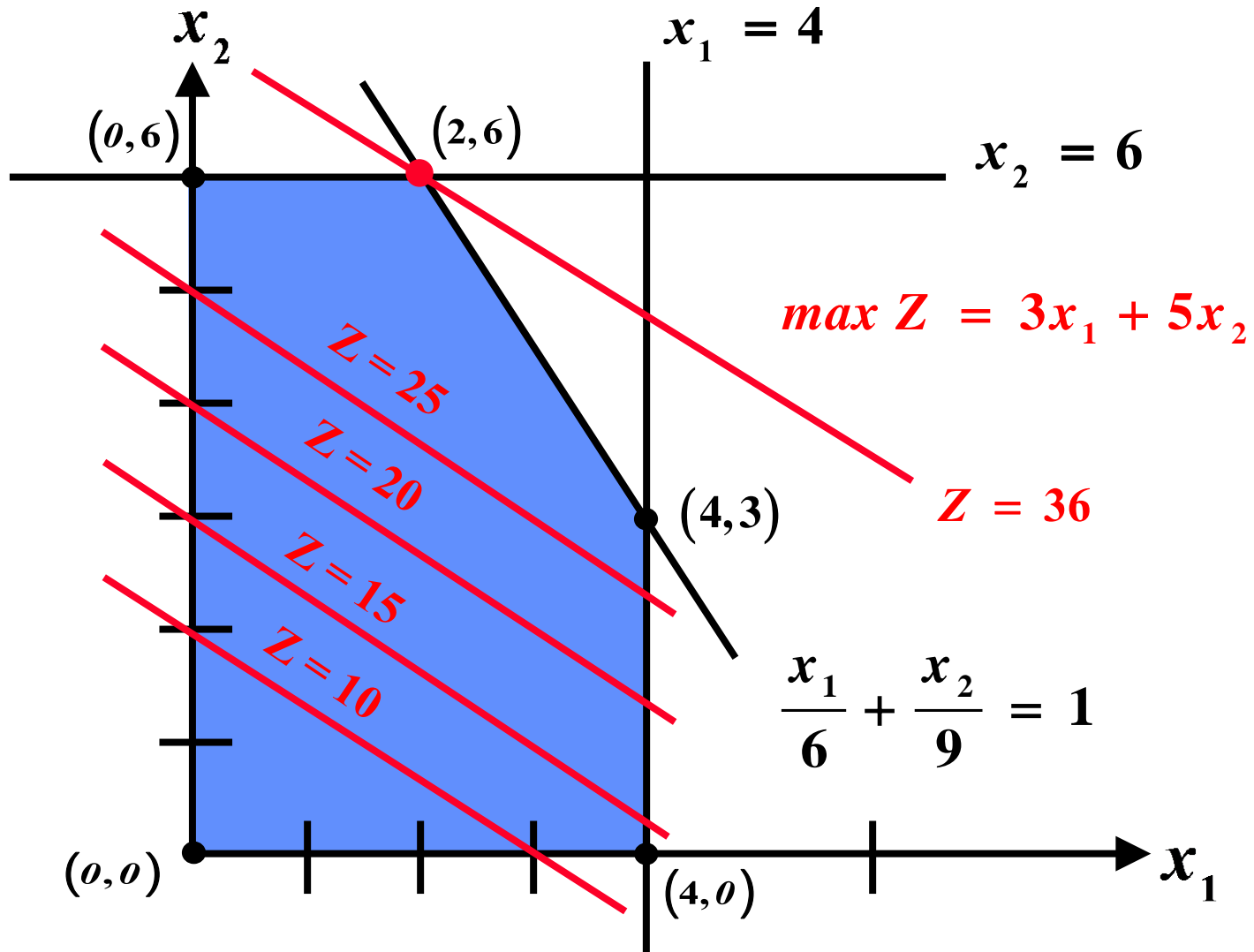
THE *FEASIBLE REGION*



FEASIBLE SOLUTION SPACE



CONTOURS OF CONSTANT Z



OPTIMAL SOLUTION

□ For this simple problem, we can *graphically* obtain the *optimal* solution

□ The *optimal* solution of this problem is:

$$x_1^* = 2 \quad \text{and} \quad x_2^* = 6$$

□ The objective value at the *optimal* solution is

$$Z^* = 3x_1^* + 5x_2^* = 36$$

LINEAR PROGRAMMING (*LP*) PROBLEM DEFINITION

A linear programming problem is an *optimization problem* with a *linear* objective function and *linear* constraints.

EXAMPLE 3: ONE-POTATO, TWO-POTATO PROBLEM

- ❑ Mr. Spud manages the *Potatoes-R-Us Co.* which processes potatoes into packages of freedom fries (F), hash browns (H) and chips (C)
- ❑ Mr. Spud can buy potatoes from two sources; each source has distinct characteristics/limits
- ❑ The problem is to determine the respective quantities Mr. Spud needs to buy from source 1 and from source 2 so as to maximize his profits

EXAMPLE 3: ONE-POTATO, TWO-POTATO PROBLEM

- ❑ The given data are summarized in the table

<i>product</i>	<i>source 1 uses (%)</i>	<i>source 2 uses (%)</i>	<i>sales limit (tons)</i>
<i>F</i>	20	30	1.8
<i>H</i>	20	10	1.2
<i>C</i>	30	30	2.4
<i>profits (\$/ton)</i>	5	6	—

- ❑ The following assumptions hold:

- 30 % waste for each source
- production may not exceed the sales limit

ANALYSIS

□ Decision variables:

$x_1 = \text{quantity purchased from source 1}$

$x_2 = \text{quantity purchased from source 2}$

□ Objective function:

$$\max Z = 5x_1 + 6x_2$$

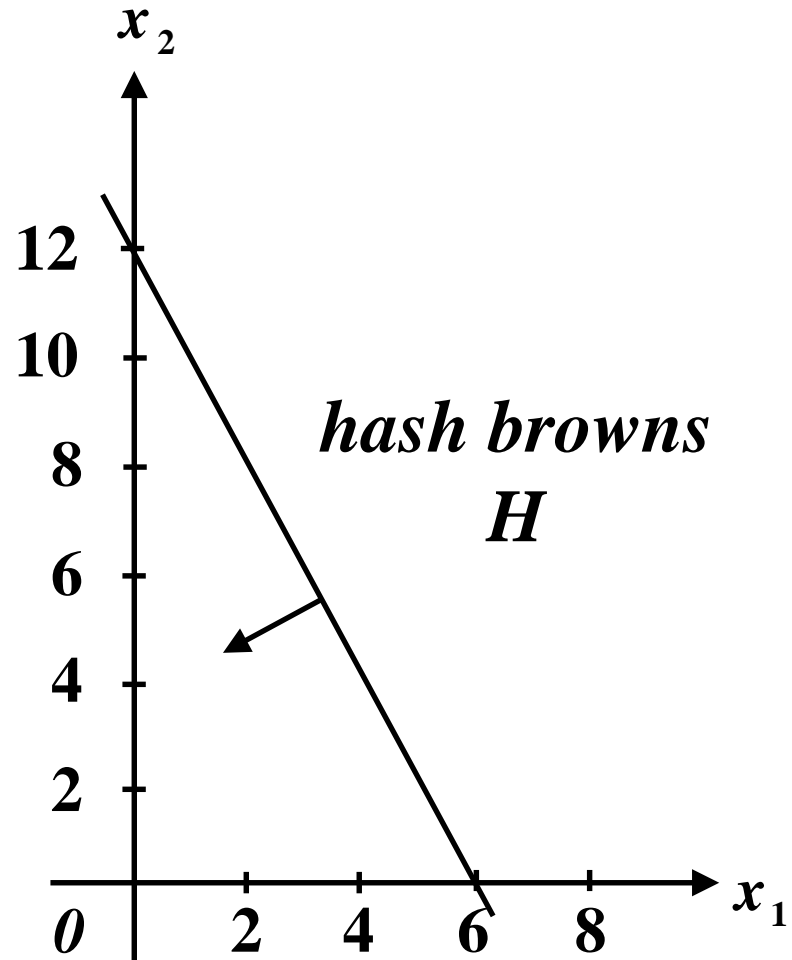
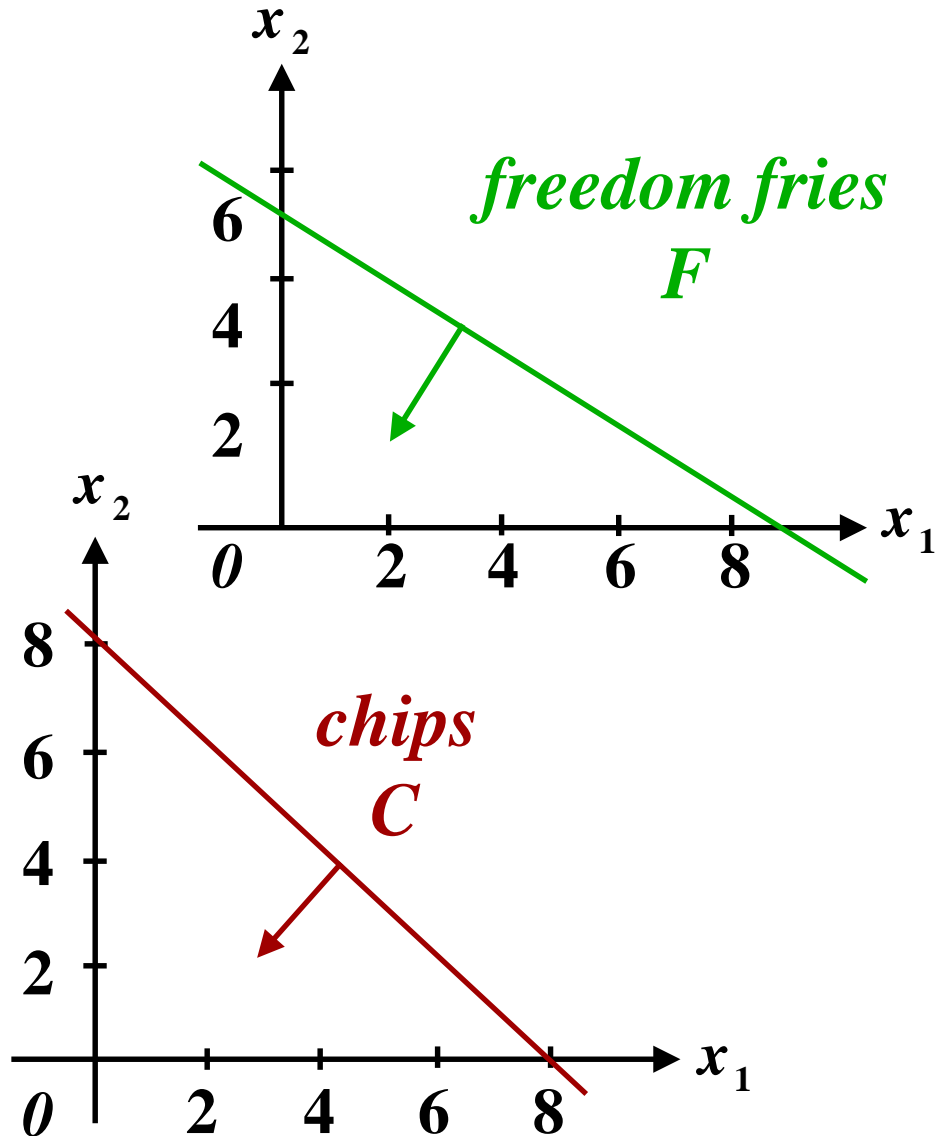
□ Constraints:

$$0.2x_1 + 0.3x_2 \leq 1.8 \quad (F)$$

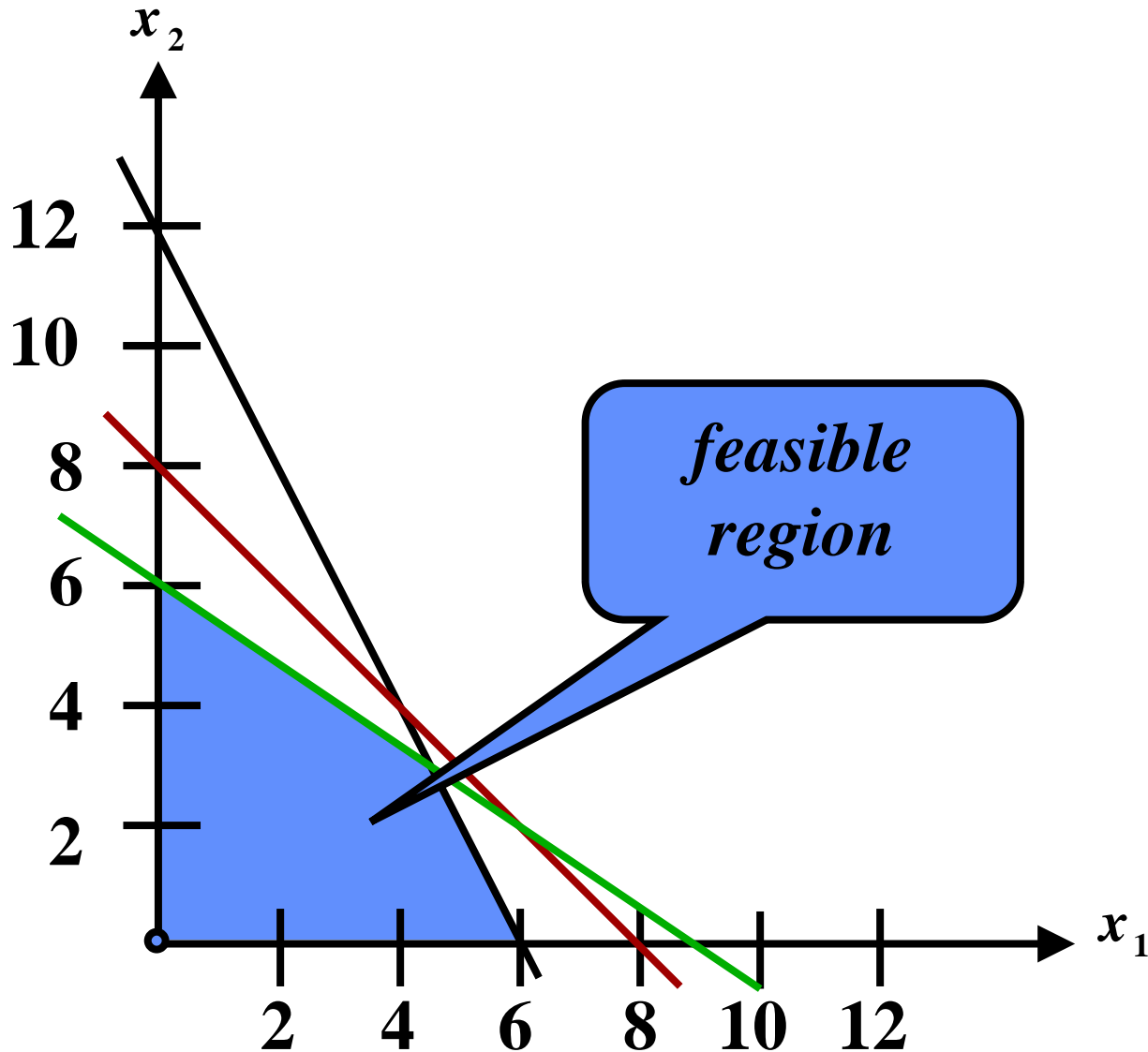
$$0.2x_1 + 0.1x_2 \leq 1.2 \quad (H) \quad x_1 \geq 0, x_2 \geq 0$$

$$0.3x_1 + 0.3x_2 \leq 2.4 \quad (C)$$

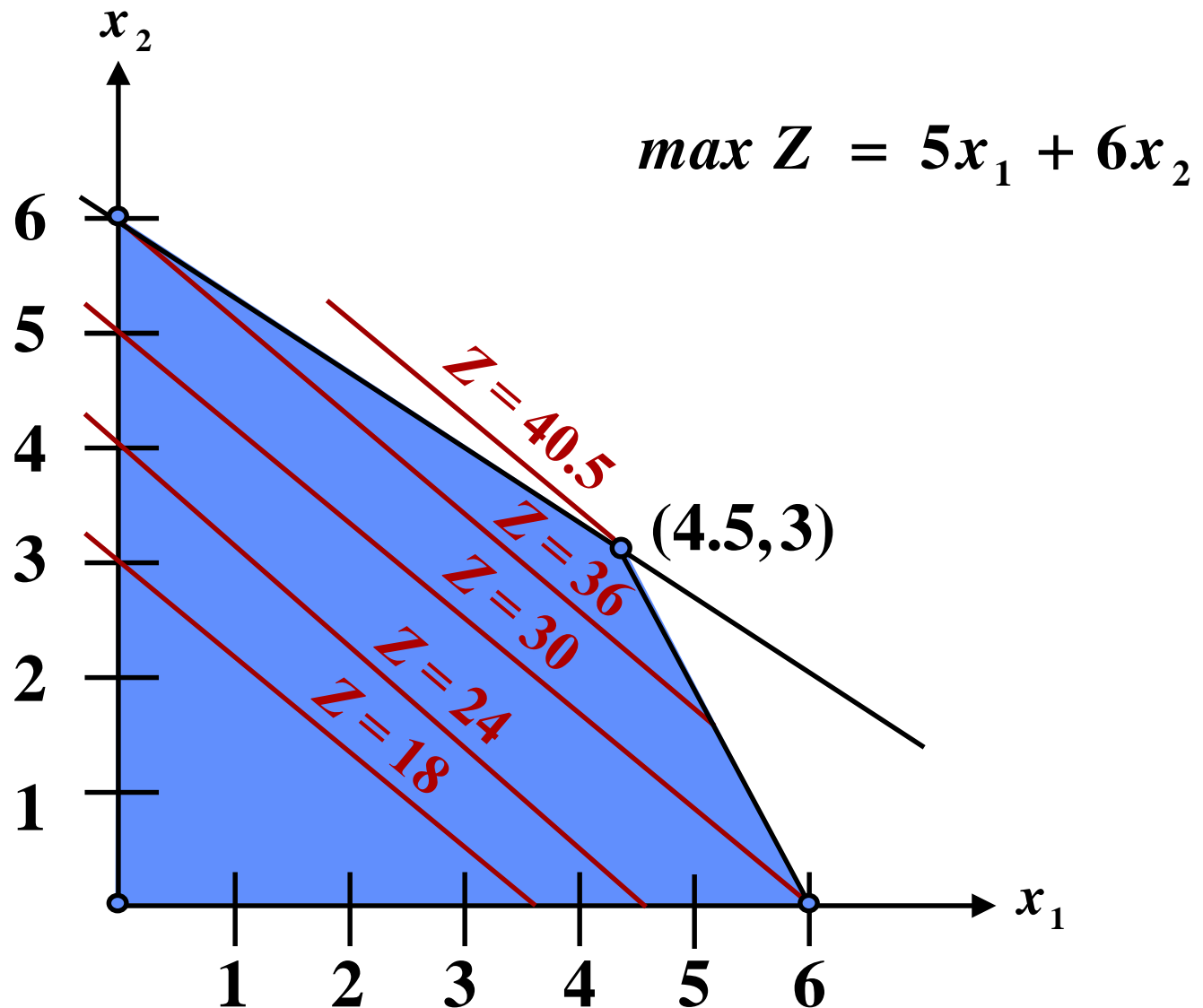
FEASIBLE REGION DETERMINATION



THE FEASIBLE REGION



EXAMPLE 3: CONTOURS OF CONSTANT Z



THE OPTIMAL SOLUTION

□ The optimal solution of this problem is:

$$x_1^* = 4.5 \qquad x_2^* = 3$$

□ The objective value at the optimal solution is:

$$Z^* = 5x_1^* + 6x_2^* = 40.5$$

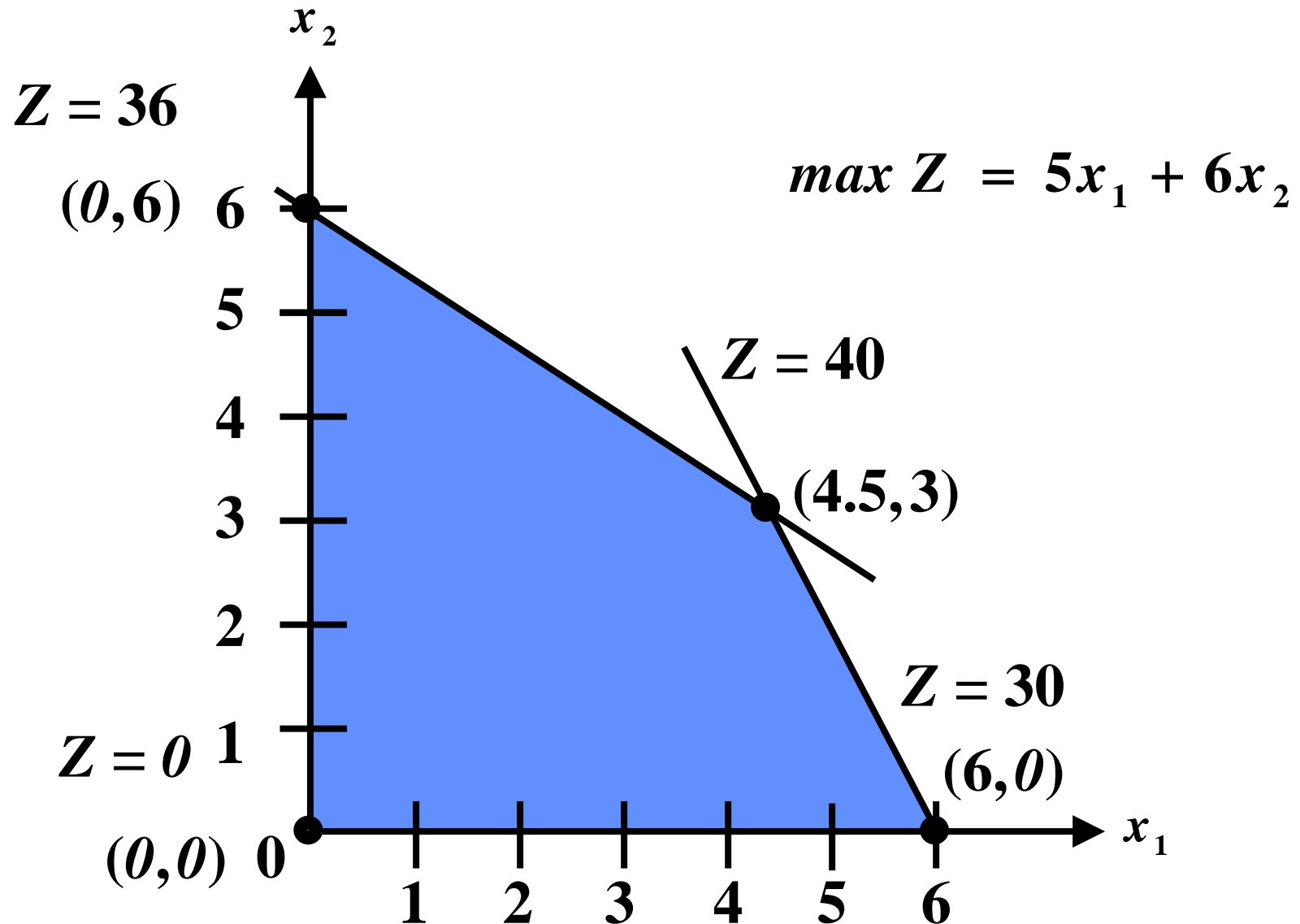
IMPORTANT OBSERVATIONS

- ❑ Constant Z lines are parallel and change monotonically along the direction normal to the contours of constant values of Z
- ❑ An *optimal* solution must be at one of the *corner points* of the feasible region: fortuitously, there are only a *finite* number of *corner points*
- ❑ If a particular *corner point* gives a better solution (in terms of its Z value) than that at every other adjacent *corner point*, then, it is an *optimal* solution

CONCEPTUAL SOLUTION PROCEDURE

- ❑ Initialization step: start at a *corner point*
- ❑ Iteration step: move to an improved *adjacent corner point* and repeat this step as many times as needed
- ❑ Stopping rule: stop when the *corner point* solution is better than that at each *adjacent corner point*
- ❑ This conceptual procedure forms the basis of the *simplex approach*

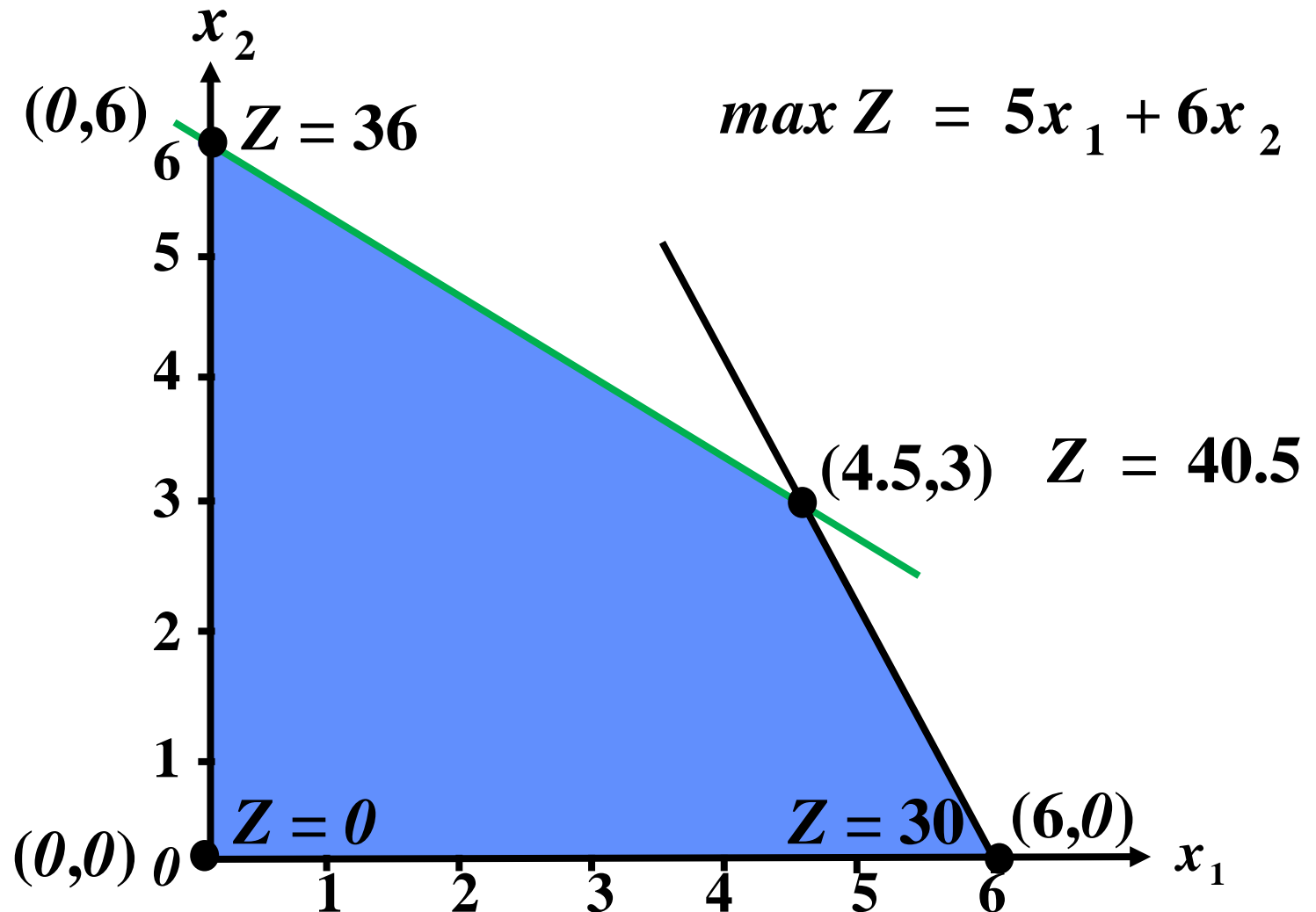
EXAMPLE 3: THE SIMPLEX APPROACH SOLUTION



EXAMPLE 3 : THE SIMPLEX APPROACH SOLUTION

<i>step</i>	x_2	x_1	Z
<i>0</i>	<i>0</i>	<i>0</i>	<i>0</i>
1	0	6	36
2	4.5	3	40.5
3	6	<i>0</i>	30

EXAMPLE 3 : THE SIMPLEX APPROACH SOLUTION



EXAMPLE 3 : THE SIMPLEX APPROACH SOLUTION

1. Start at $(0,0)$ with $Z(0,0) = 0$
2. (i) Move from $(0,0)$ to $(0,6)$, $Z(0,6) = 36$
(ii) Move from $(0,6)$ to $(4.5,3)$; compute $Z(4.5,3) = 40.5$
3. Compare the objective at $(4.5,3)$ to values at $(6,0)$ and at $(0,6)$:

$$Z(4.5,3) \geq Z(6,0)$$

$$Z(4.5,3) \geq Z(0,6)$$

therefore, $(4.5,3)$ is *optimal*

REVIEW

- ❑ Key requirements of a programming problem:
 - to make a decision, we must define the *decision variables*
 - to achieve the specified objective, we must express mathematically the *objective function*
 - to ensure *feasibility*, the decision variables must satisfy each *mathematically formulated constraint*

REVIEW

❑ Key attributes of an *LP*

- the objective function is *linear*
- the constraints are *linear*

❑ Basic steps in formulating a programming problem

- definition of decision variables
- statement of the objective function
- formulation of the constraints

REVIEW

- ❑ Words of caution: care is required with units and attention is needed to not ignore the *implicit constraints*, such as nonnegativity, and the common sense requirements in an *LP* formulation
- ❑ Graphical solution approach for two–variable problems
 - feasible region determination
 - contours of constant Z
 - identification of the vertex with optimal Z^*

EXAMPLE 4 : QUALITY CONTROL INSPECTION OF GOODS PRODUCED

- ❑ There are 8 grade 1 and 10 grade 2 inspectors available for *QC* inspection; at least 1,800 pieces must be inspected in each 8-*hour* day
- ❑ Problem data are summarized below:

<i>grade level</i>	<i>speed (unit/h)</i>	<i>accuracy (%)</i>	<i>wages (\$/h)</i>
1	25	98	4
2	15	95	3

EXAMPLE 4 : INSPECTION OF GOODS PRODUCED

❑ Each error costs \$ 2

❑ The problem is to determine the *optimal*

assignment of inspectors, i.e., the number of

inspectors of grade 1 and that of grade 2 to result

in the least-cost *QC* inspection effort

EXAMPLE 4 : FORMULATION

□ Definition of decision variables:

x_1 = number of grade 1 inspectors assigned

x_2 = number of grade 2 inspectors assigned

□ Objective function

○ optimal assignment: minimum costs

○ costs = wages + errors

EXAMPLE 4 : FORMULATION

- each grade 1 inspector costs:

$$4 + 2 (25) (0.02) = 5 \$ / hr$$

- each grade 2 inspector costs:

$$3 + 2 (15) (0.05) = 4.5 \$ / hr$$

- total daily inspection costs in \$ are

$$Z = 8[5x_1 + 4.5x_2] = 40x_1 + 36x_2 \quad (\$)$$

EXAMPLE 4 : FORMULATION

□ Constraints:

○ job completion:

$$8(25)x_1 + 8(15)x_2 \geq 1,800$$

$$\Leftrightarrow 200x_1 + 120x_2 \geq 1,800$$

$$\Leftrightarrow 5x_1 + 3x_2 \geq 45$$

○ availability limit:

$$x_1 \leq 8$$

$$x_2 \leq 10$$

○ nonnegativity:

$$x_1 \geq 0, x_2 \geq 0$$

EXAMPLE 4 : PROBLEM STATEMENT SUMMARY

□ Decision variables:

x_1 = number of grade 1 inspectors assigned

x_2 = number of grade 2 inspectors assigned

□ Objective function:

$$\min Z = 40x_1 + 36x_2$$

□ Constraints:

$$5x_1 + 3x_2 \geq 45$$

$$x_1 \leq 8$$

$$x_2 \leq 10$$

$$x_1 \geq 0, x_2 \geq 0$$

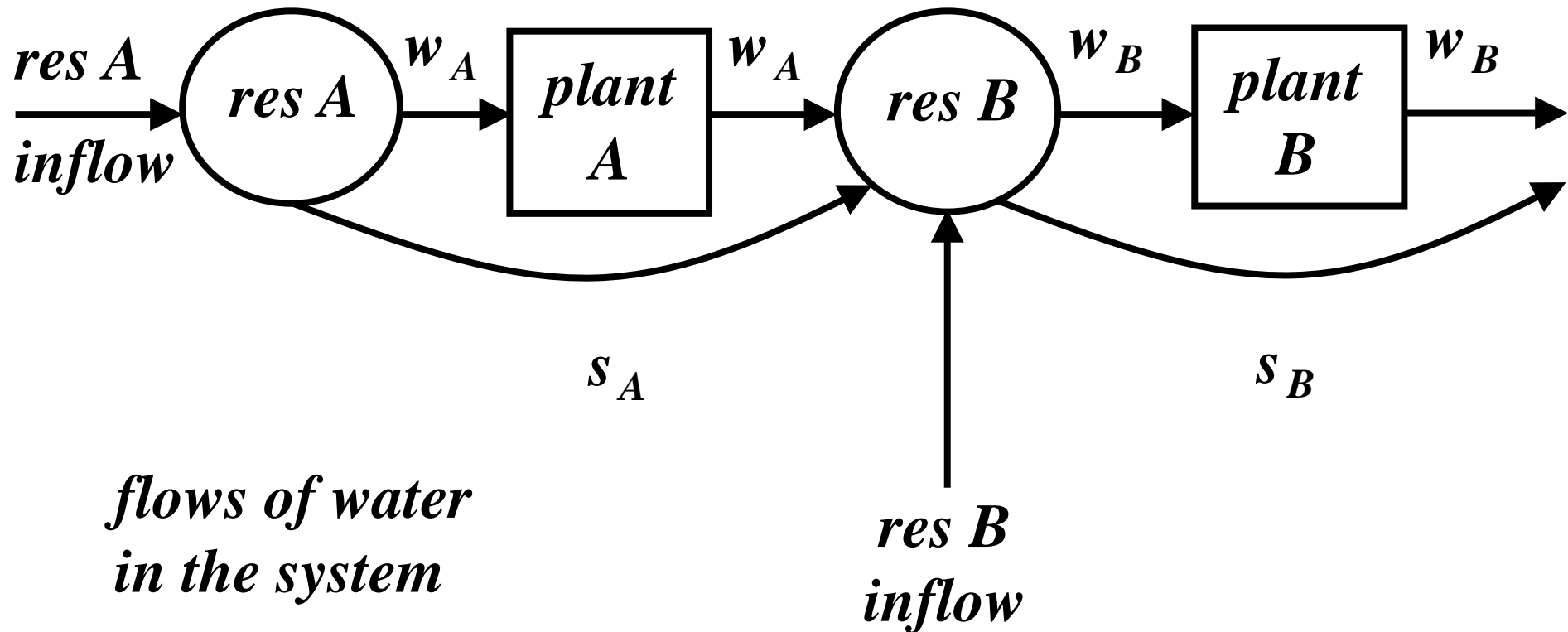
MULTI – PERIOD SCHEDULING

- ☐ **More than one period is involved**
- ☐ **The result of each period affects the initial conditions for the next period and therefore the solution**
- ☐ **We need to define variables to take into account the initial conditions in addition to the decision variables of the problem**

EXAMPLE 5 : HYDROELECTRIC POWER SYSTEM OPERATIONS

- We consider a single operator of a system consisting of two water reservoirs with a hydroelectric plant attached to each reservoir**
- We schedule the two power plant operations over a two-period horizon**
- We are interested in a plan to maximize the total revenues of the system operator**

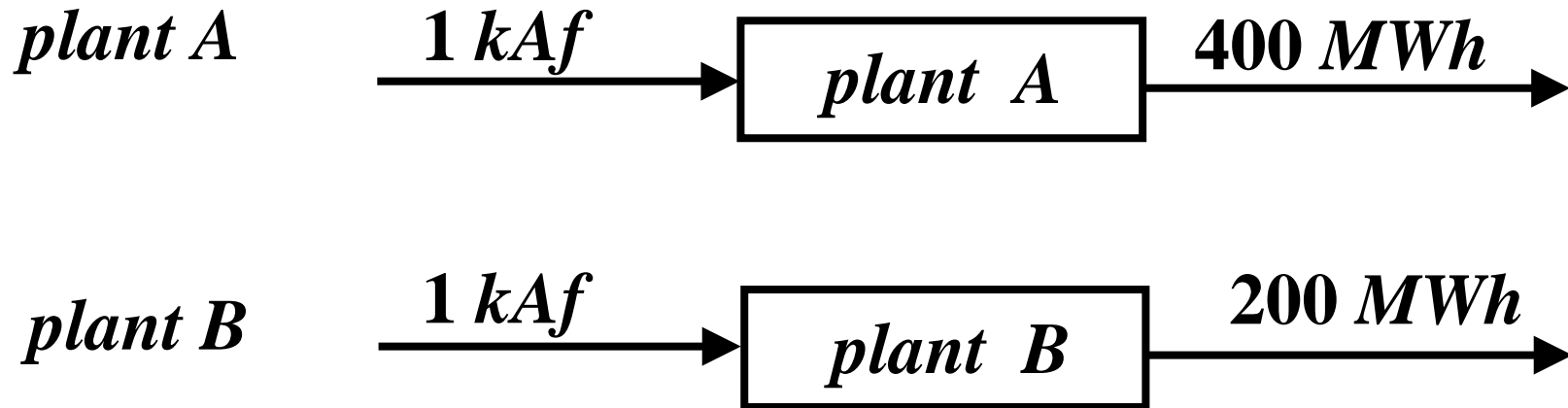
EXAMPLE 5 : HYDROELECTRIC POWER SYSTEM OPERATIONS



EXAMPLE 5 : kAf RESERVOIR DATA

<i>parameter</i>	<i>reservoir A</i>	<i>reservoir B</i>
<i>maximum capacity</i>	2,000	1,500
<i>predicted inflow in period 1</i>	200	40
<i>predicted inflow in period 2</i>	130	15
<i>minimum allowable level</i>	1,200	800
<i>level at start of period 1</i>	1,900	850

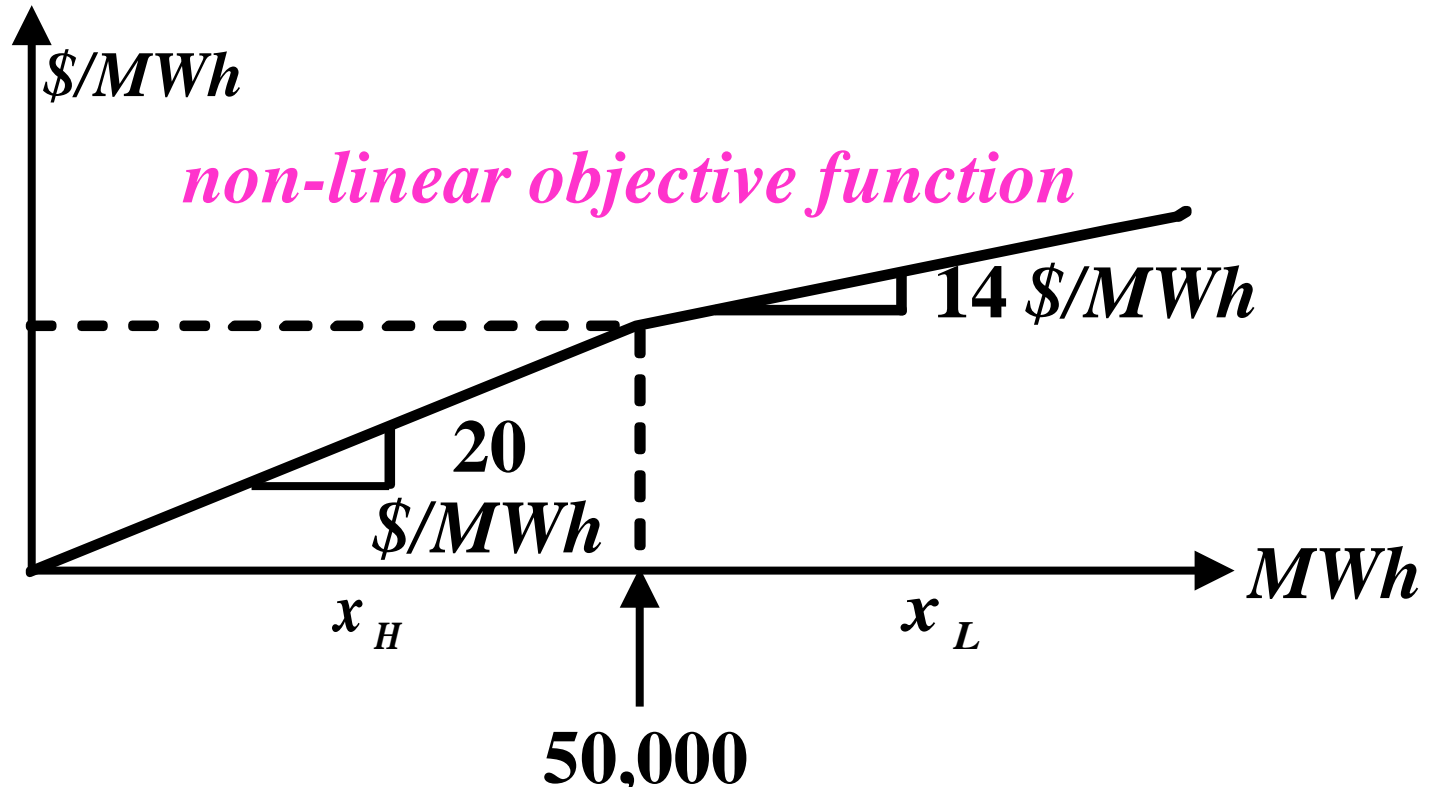
EXAMPLE 5 : SYSTEM CHARACTERISTICS



<i>reservoir</i>	<i>max kAf for generation per period</i>
A	150
B	87.5

EXAMPLE 5 : SYSTEM CHARACTERISTICS

- ❑ Two-tier price for the MWh demand in each period
 - up to 50,000 MWh can be sold @ 20 \$ / MWh
 - all additional MWh are sold @ 14 \$ / MWh



EXAMPLE 5 : DECISION VARIABLES

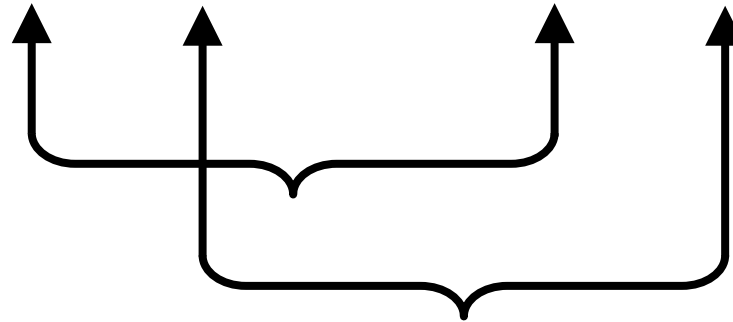
<i>variable</i>	<i>quantity denoted</i>	<i>units</i>
x_H^i	<i>energy sold at 20 \$/MWh</i>	<i>MWh</i>
x_L^i	<i>energy sold at 14 \$/MWh</i>	<i>MWh</i>
w_A^i	<i>plant A water supply for generation</i>	<i>kAf</i>
w_B^i	<i>plant B water supply for generation</i>	<i>kAf</i>
s_A^i	<i>reservoir A spill</i>	<i>kAf</i>
s_B^i	<i>reservoir B spill</i>	<i>kAf</i>
r_A^i	<i>reservoir A end of period i level</i>	<i>kAf</i>
r_B^i	<i>reservoir B end of period i level</i>	<i>kAf</i>

superscript i denotes period i , $i = 1, 2$

EXAMPLE 5 : OBJECTIVE FUNCTION

maximize total revenues from sales

$$\max \quad Z = 20(x_H^1 + x_H^2) + 14(x_L^1 + x_L^2)$$



4 of the 16 decision variables
2 for each period

units of Z are in \$

EXAMPLE 5 : CONSTRAINTS

□ Period 1 constraints

○ energy conservation in a lossless system

- total generation $400w_A^1 + 200w_B^1$ (MWh)
- total sales $x_H^1 + x_L^1$ (MWh)
- losses are neglected and so

$$x_H^1 + x_L^1 = 400w_A^1 + 200w_B^1$$

○ maximum available capacity limits

$$w_A^1 \leq 150$$

$$w_B^1 \leq 87.5$$

EXAMPLE 5 : CONSTRAINTS

○ reservoir conservation of flow relations

- reservoir *A*:

$$w_A^1 + s_A^1 + r_A^1 = 1,900 + 200 = 2,100 \text{ (kAf)}$$

*res. level at
e.o.p. 1*

*res. level at
e.o.p. 0*

*predicted
inflow*

- reservoir *B*:

$$w_B^1 + s_B^1 + r_B^1 = 850 + 40 + w_A^1 + s_A^1 \text{ (kAf)}$$

EXAMPLE 5 : CONSTRAINTS

○ limitations on reservoir variables

- reservoir A :

$$1,200 \leq r_A^1 \leq 2,000 \quad (kAf)$$

- reservoir B :

$$800 \leq r_B^1 \leq 1,500 \quad (kAf)$$

○ sales constraint

$$x_H^1 \leq 50,000 \quad (kAf)$$

EXAMPLE 5 : CONSTRAINTS

□ Period 2 constraints

○ energy conservation in a lossless system

- total generation $400w_A^2 + 200w_B^2$ (MWh)

- total sales $x_H^2 + x_L^2$ (MWh)

- losses are neglected and so

$$x_H^2 + x_L^2 = 400w_A^2 + 200w_B^2$$

○ maximum available capacity limits

$$w_A^2 \leq 150$$

$$w_B^2 \leq 87.5$$

EXAMPLE 5 : CONSTRAINTS

○ reservoir conservation of flow relations

- reservoir A:

$$w_A^2 + s_A^2 + r_A^2 = r_A^1 + 130 \quad (kAf)$$

*res. level at
e.o.p. 2*

*res. level at
e.o.p. 1*

*predicted
inflow*

- reservoir B:

$$w_B^2 + s_B^2 + r_B^2 = r_B^1 + 15 + w_A^2 + s_A^2 \quad (kAf)$$

EXAMPLE 5 : CONSTRAINTS

○ limitations on reservoir variables

• reservoir A :

$$1,200 \leq r_A^2 \leq 2,000 \quad (kAf)$$

• reservoir B :

$$800 \leq r_B^2 \leq 1,500 \quad (kAf)$$

○ sales constraint

$$x_H^2 \leq 50,000 \quad (kAf)$$

EXAMPLE 5 : PROBLEM STATEMENT

□ 16 decision variables:

$$x_H^i, x_L^i, w_A^i, w_B^i, s_A^i, s_B^i, r_A^i, r_B^i, \quad i = 1, 2$$

□ Objective function:

$$\max \quad Z = 20(x_H^1 + x_H^2) + 14(x_L^1 + x_L^2)$$

□ Constraints:

- 20 constraints for the periods 1 and 2
- non-negativity constraints on all variables

EXAMPLE 6 : DISHWASHER AND WASHING MACHINE PROBLEM

- ❑ The *Appliance Co.* manufactures dishwashers and washing machines
- ❑ The sales targets for next four quarters are:

<i>product</i>	<i>variable</i>	<i>quarter t</i>			
		1	2	3	4
<i>dishwasher</i>	D_t	2,000	1,300	3,000	1,000
<i>washing machine</i>	W_t	1,200	1,500	1,000	1,400

EXAMPLE 6 : QUARTERLY COST COMPONENTS

<i>cost component</i>		<i>parameter</i>	<i>quarter t costs (\$/unit)</i>			
			1	2	3	4
<i>manufacturing (\$/unit)</i>	<i>dishwasher</i>	c_t	125	130	125	126
	<i>washing machine</i>	v_t	90	100	95	95
<i>storage (\$/unit)</i>	<i>dishwasher</i>	j_t	5.0	4.5	4.5	4.0
	<i>washing machine</i>	k_t	4.3	3.8	3.8	3.3
<i>hourly labor (\$ /hour)</i>		p_t	6.0	6.0	6.8	6.8

EXAMPLE 6 : CONSTRAINTS

- ❑ Each dishwasher (washing machine) requires 1.5
(2) *hours* of labor
- ❑ The labor hours in each quarter cannot grow or decrease by more than 10 %; there are 5,000 *h* of labor in the quarter preceding the first quarter
- ❑ At the start of the first quarter, there are 750 dishwashers and 50 washing machines in storage

EXAMPLE 6 : THE PROBLEM

**How to schedule the production in each of the
four quarters so as to minimize the costs while
meeting the sales targets?**

EXAMPLE 6 : QUARTER t DECISION VARIABLES

<i>symbol</i>	<i>variable</i>
d_t	<i>number of dishwashers produced</i>
w_t	<i>number of washing machines produced</i>
r_t	<i>final inventory of dishwashers</i>
s_t	<i>final inventory of washing machines</i>
h_t	<i>available labor hours during Q_t</i>
$t = 1, 2, 3, 4$	

EXAMPLE 6 : OBJECTIVE FUNCTION

minimize the *total* costs for the four quarters

*manufacturing
costs*

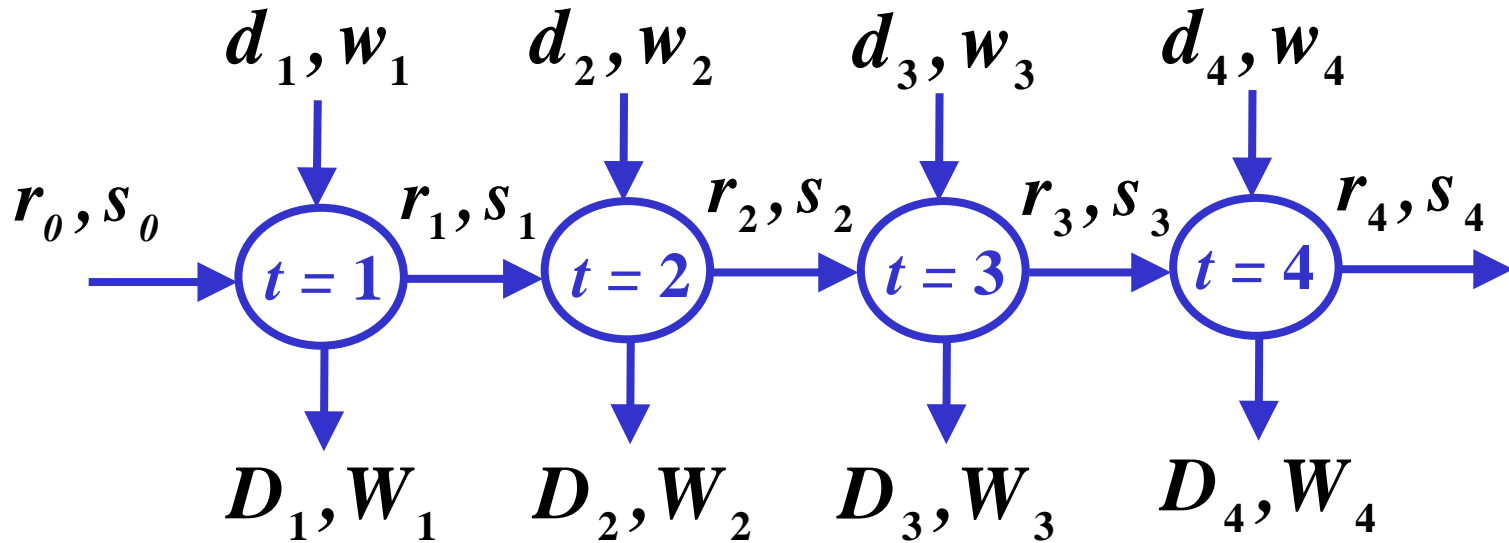
*storage
costs*

*labor
costs*

$$\begin{aligned} \min Z = & \underbrace{c_1 d_1 + v_1 w_1}_{\text{manufacturing costs}} + \underbrace{j_1 r_1 + k_1 s_1}_{\text{storage costs}} + \underbrace{p_1 h_1}_{\text{labor costs}} \leftarrow \text{quarter 1} \\ & + c_2 d_2 + v_2 w_2 + j_2 r_2 + k_2 s_2 + p_2 h_2 \leftarrow \text{quarter 2} \\ & + c_3 d_3 + v_3 w_3 + j_3 r_3 + k_3 s_3 + p_3 h_3 \leftarrow \text{quarter 3} \\ & + c_4 d_4 + v_4 w_4 + j_4 r_4 + k_4 s_4 + p_4 h_4 \leftarrow \text{quarter 4} \end{aligned}$$

EXAMPLE 6 : CONSTRAINTS

□ Quarterly flow balance relations:



$$\begin{cases} r_{t-1} + d_t - r_t = D_t \\ s_{t-1} + w_t - s_t = W_t \end{cases} \quad t = 1, 2, 3, 4$$

EXAMPLE 6 : CONSTRAINTS

□ Quarterly labor constraints

$$\begin{cases} 1.5d_t + 2w_t - h_t \leq 0 \\ 0.9h_{t-1} \leq h_t \leq 1.1h_{t-1} \end{cases} \quad t = 1, 2, 3, 4$$

$$h_0 = 5,000$$

EXAMPLE 6 : PROBLEM STATEMENT

d_1	w_1	r_1	s_1	h_1	d_2	w_2	r_2	s_2	h_2	d_3	w_3	r_3	s_3	h_3	d_4	w_4	r_4	s_4	h_4	
1		-1																		= 1250
	1		-1																	= 1150
1.5	2			-1																≤ 0
				1																≥ 4500
				1																≤ 5500
		1			1		-1													= 1300
			1			1		-1												= 1500
					1.5	2			-1											≤ 0
				-0.9					1											≥ 0
				-1.1					1											≤ 0
							1			1		-1								= 3000
								1			1		-1							= 1000
										1.5	2			-1						≤ 0
									-0.9					1						≥ 0
									-1.1					1						≤ 0
												1			1		-1			= 1000
													1			1		-1		= 1400
															1.5	2			-1	≤ 0
														-0.9					1	≥ 0
														-1.1					1	≤ 0
125	90	5.0	4.3	6.0	130	100	4.5	3.8	6.0	125	95	4.5	3.8	6.8	126	95	4.0	3.3	6.8	<i>minimize</i>

LINEAR PROGRAMMING PROBLEM

$$\max (\min) \quad Z = c_1 x_1 + \dots + c_n x_n$$

s.t.

$$a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n = b_1$$

$$a_{21} x_1 + a_{22} x_2 + \dots + a_{2n} x_n = b_2$$

$$\vdots$$
$$\vdots$$

$$a_{m1} x_1 + a_{m2} x_2 + \dots + a_{mn} x_n = b_m$$

$$x_1 \geq 0, x_2 \geq 0, \dots, x_n \geq 0$$

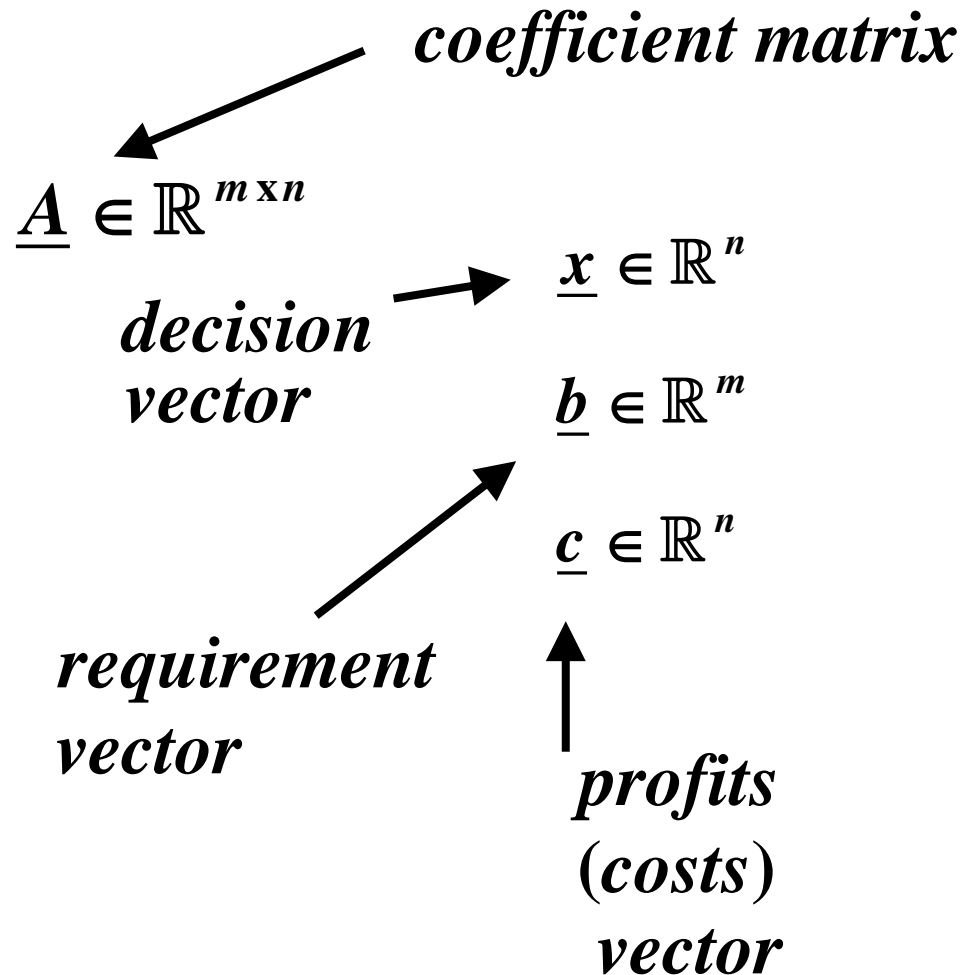
$$b_1 \geq 0, b_2 \geq 0, \dots, b_m \geq 0$$

STANDARD FORM OF *LP* (*SFLP*)

$$\max(\min) Z = \underline{c}^T \underline{x}$$

$$\underline{A} \underline{x} = \underline{b}$$

$$\underline{x} \geq \underline{0}$$



CONVERSION OF *LP* INTO *SFLP*

□ An inequality may be converted into an equality by defining an additional nonnegative *slack* variable

○ $x_{slack} \geq 0$

○ replace the given *inequality* $\leq b$ by

$$inequality + x_{slack} = b$$

○ replace the given *inequality* $\geq b$ by

$$inequality - x_{slack} = b$$

CONVERSION OF LP INTO $SFLP$

- An unsigned variable x_u is one whose sign is *not* specified
- x_u may be converted into two signed variables x_+ and x_- with

$$x_+ = \begin{cases} x_u & x_u \geq 0 \\ 0 & x_u < 0 \end{cases} \quad x_- = \begin{cases} 0 & x_u \geq 0 \\ -x_u & x_u < 0 \end{cases}$$

so that x_u is replaced by

$$x_u = x_+ - x_-$$

SFLP CHARACTERISTICS

- \underline{x} is feasible if and only if $\underline{x} \geq \underline{0}$ and $\underline{A}\underline{x} = \underline{b}$
- $S = \{\underline{x} \mid \underline{A}\underline{x} = \underline{b}, \underline{x} \geq \underline{0}\}$ is the feasible region
- $S = \emptyset \Rightarrow LP$ is infeasible
- \underline{x}^* is optimal $\Rightarrow \underline{c}^T \underline{x}^* \geq \underline{c}^T \underline{x}, \underline{x} \in S$
- \underline{x}^* may be unique, or may have multiple values
- \underline{x}^* may be unbounded